

# PASAN: Detecting Peripheral Access Concurrency Bugs within Bare-metal Embedded Applications

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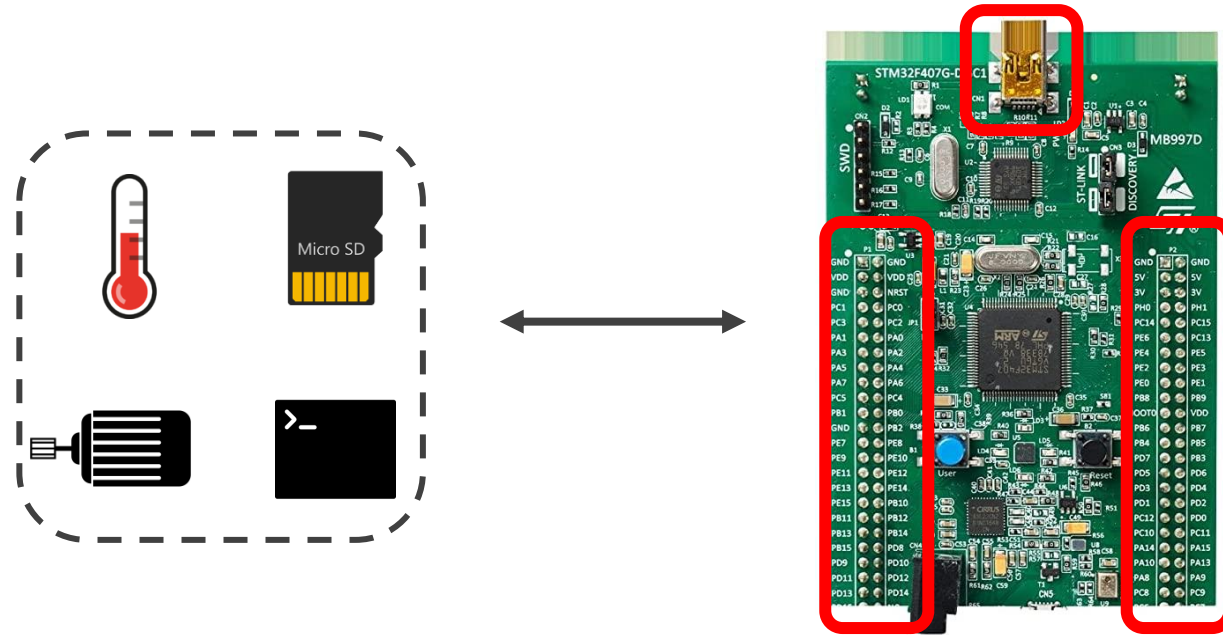
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Kyungtae Kim<sup>1</sup>, Chunghwan Kim<sup>4</sup>, Dongyan Xu<sup>1</sup>, Dave (Jing) Tian<sup>1</sup>

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# Background

- Peripheral Devices in Embedded Systems



- Attached to MMIO ports (e.g., UART, USB, SPI, etc)

# Concurrency Bugs for "Transactions"

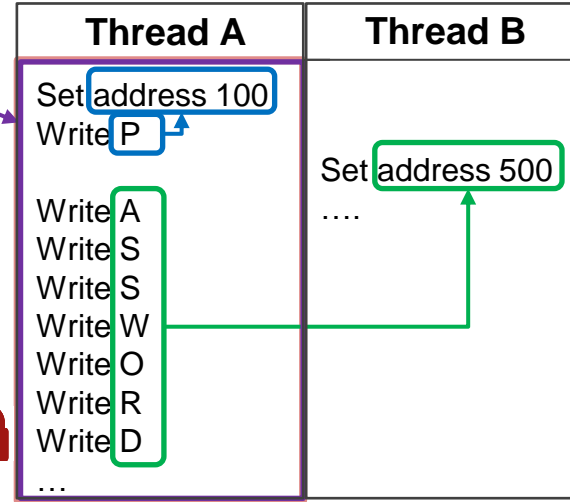
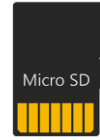
```
lock(&lock_var);  
count++; // global variable  
unlock(&lock_var);
```



Thread A	Thread B
Read <b>count</b>	Read <b>count</b>
Write <b>count + 1</b>	Write <b>count + 1</b>

## Example of Transaction

Data update on SD card



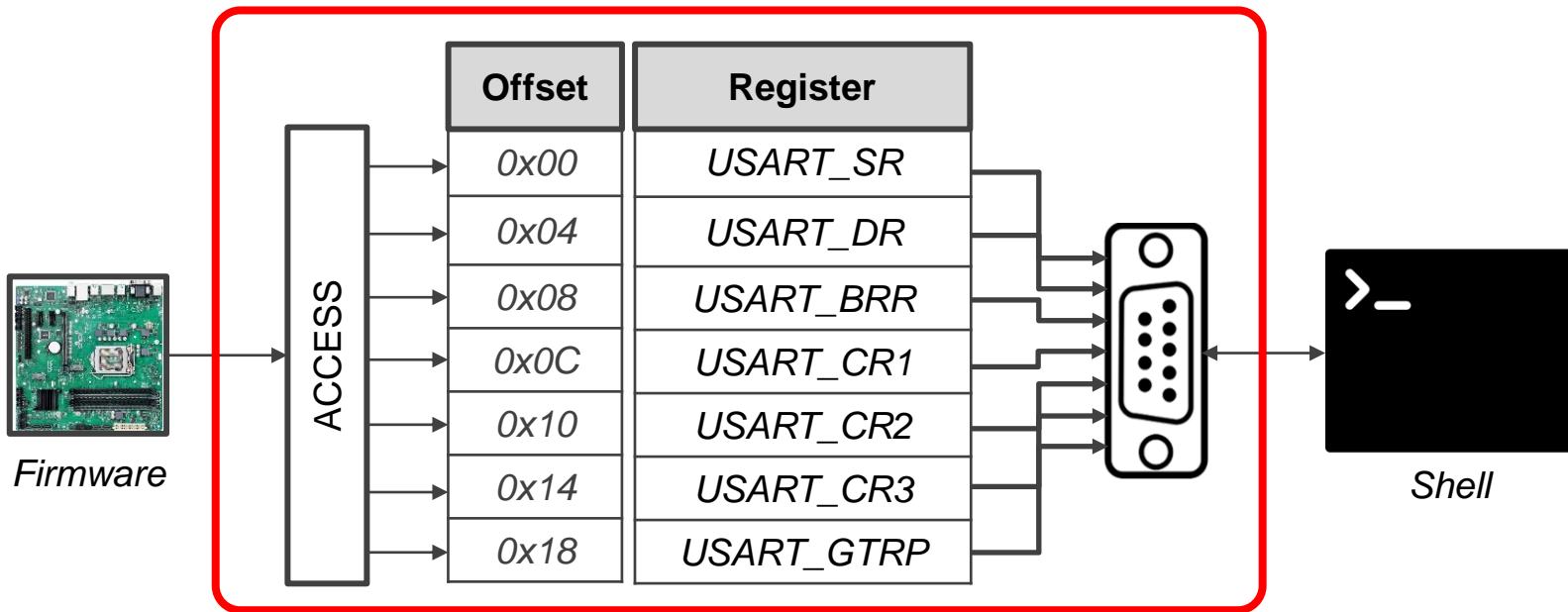
- "count" → Shared resource
- Atomic update necessary

- Peripheral devices → Shared resources
- Two aspects of atomicity guarantee
  - Spatial (address-range) atomicity
  - Temporal (concurrent execution) atomicity

Otherwise, data corruption happens (e.g., file content, sensor value, etc)

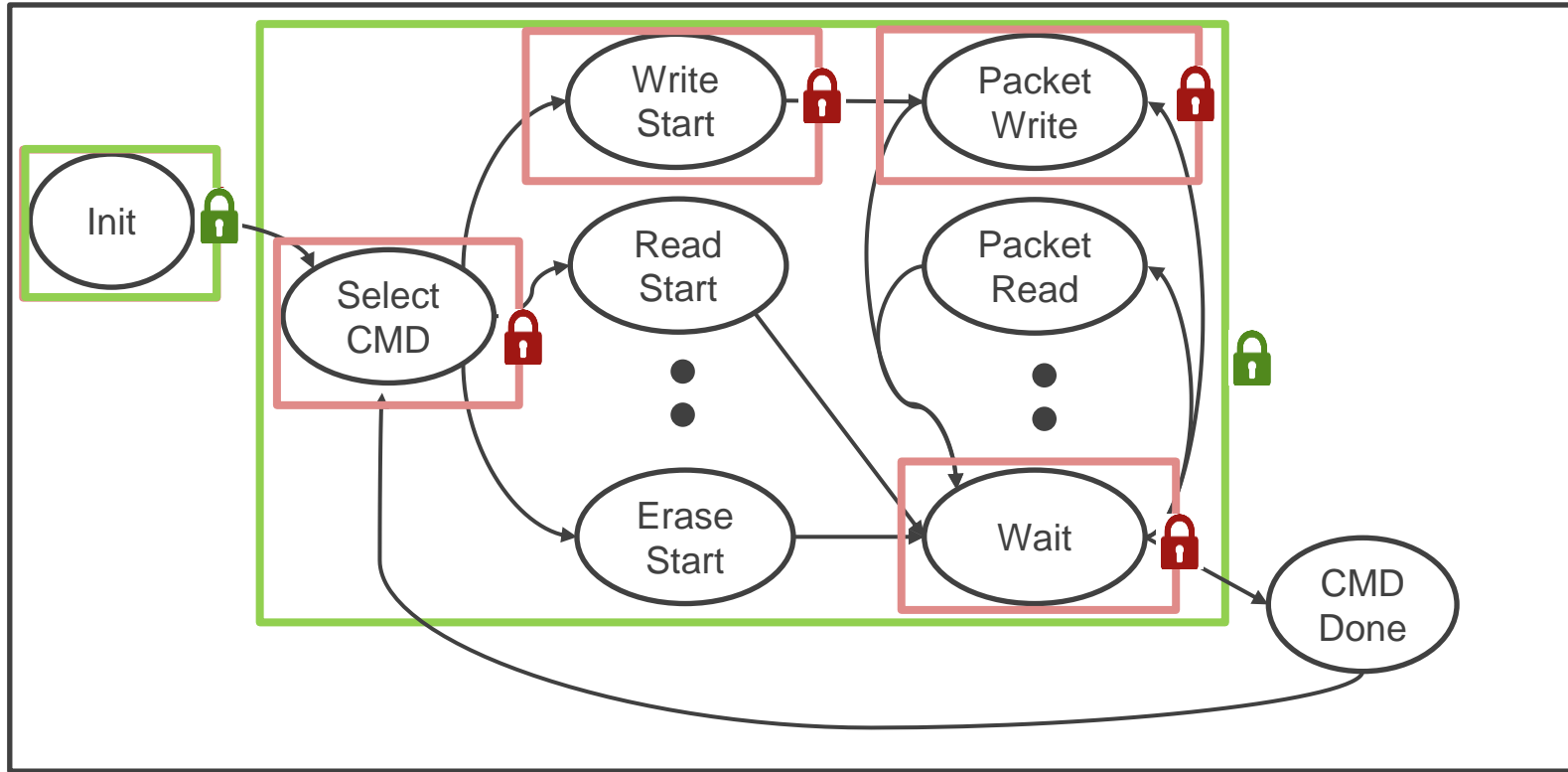
# Background: Spatial Atomicity

UART Start (Base) Address:  $0x40004400$   
UART End Address:  $0x4000441C$  } = One MMIO Region  
i.e., Spatial Atomicity



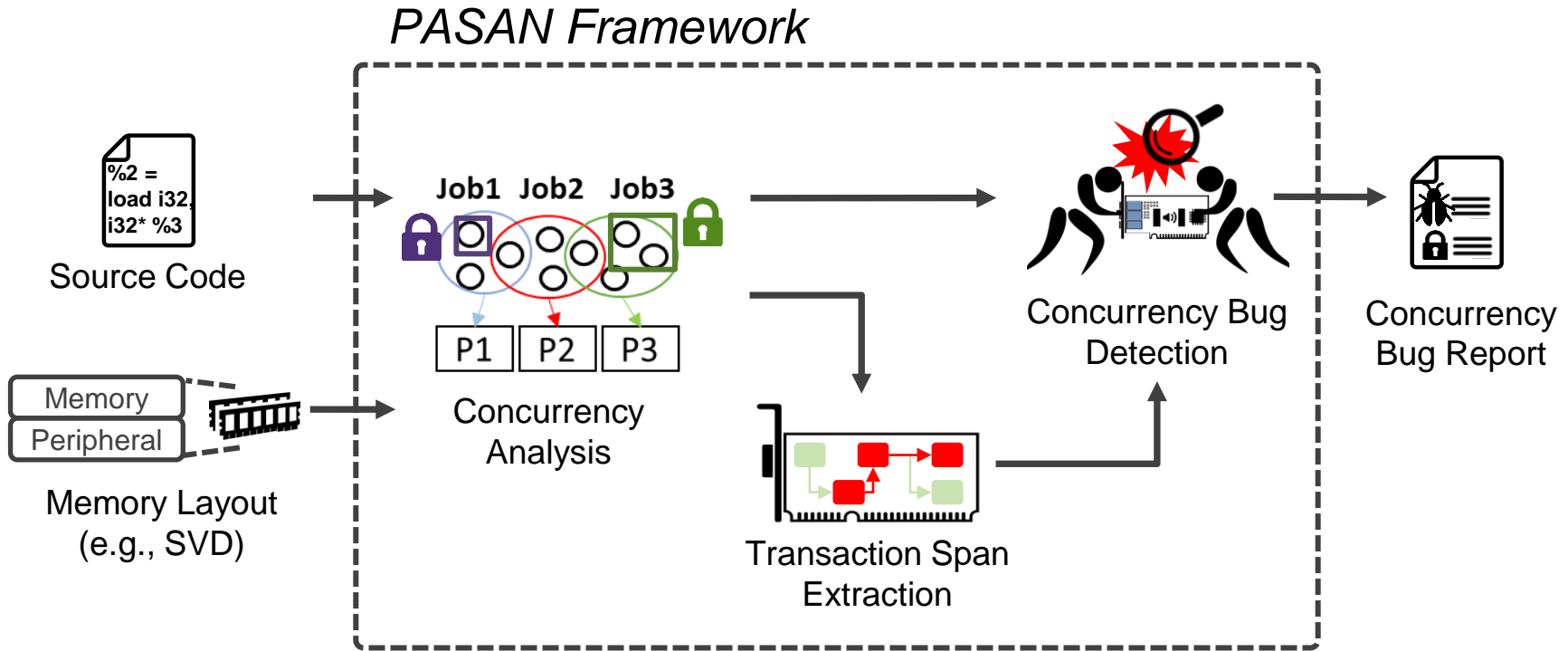
# Background: Temporal Atomicity

Simplified  
SD Card  
Controller



**“Temporal atomicity” should be guaranteed**

# PASAN Workflow



# MMIO and Lock Alias Identification

Start Address: 0x40047400  
End Address: 0x4004741C

```
void spi_write(){  
    ...  
    lock(&lock1);  
    *0x40047400 = 0x10;  
    unlock(&lock1);  
    ....  
    lock(&lock2);  
    *0x40047404 = 0x20;  
    unlock(&lock2);  
    ....  
}
```

Within Identical  
MMIO Region

```
uint32_t *mmio_ptr = 0x40047400;  
mutex_t lock;  
  
void spi_write1(){  
    ...  
    lock(&lock);  
    *(mmio_ptr) = 0x10;  
    unlock(&lock);  
    ....  
}  
  
void spi_write2(){  
    lock(&lock);  
    *(mmio_ptr + 4) = 0x20;  
    unlock(&lock);  
    ....  
}
```

# Lock Span Identification

```
void spi_write(){  
    ...  
    lock(&lock1);  
    *0x40047400 = 0x20;  
    unlock(&lock1);  
    ...  
    lock(&lock2);  
    *0x40047404 = 0x10;  
    unlock(&lock2);  
    ...  
}
```

Lock Span 1

Lock Span 2

```
void spi_write(){  
    lock(&lock);  
    spi_write1();  
    ...  
    spi_write2();  
    unlock(&lock);  
    ...  
}  
void spi_write1(){  
    ...  
    *0x40047400 = 0x10;  
    ...  
}  
void spi_write2(){  
    ...  
    *0x40047404 = 0x20;  
    ...  
}
```

Lock Span

Lock Span

Call Context -Aware



# Transaction Span Analysis

```
void spi_cmd(){
```

```
...  
lock(&lock1);
```

```
*0x40047400 = 0x10;
```

```
....  
9 instructions
```

```
....  
*0x40047404 = 0x20;
```

```
unlock(&lock2);
```

```
....  
}
```



- Distance-based Span Detection
  - Count # of instructions between MMIO access instructions
  - e.g., 10 instructions < threshold
- One exception example we should consider
  - “Wait state” after a job request
    - e.g., sleep or waiting for job completion (i.e., interrupt)
  - What if “wait state” is interrupted by another MMIO access?
    - e.g., Ongoing job can fail

# Concurrency Bug Detection

```
void spi_write(){
```

```
lock(&lock);;
```

```
*0x40047400 = 0x20;
```

```
unlock(&lock);;
```

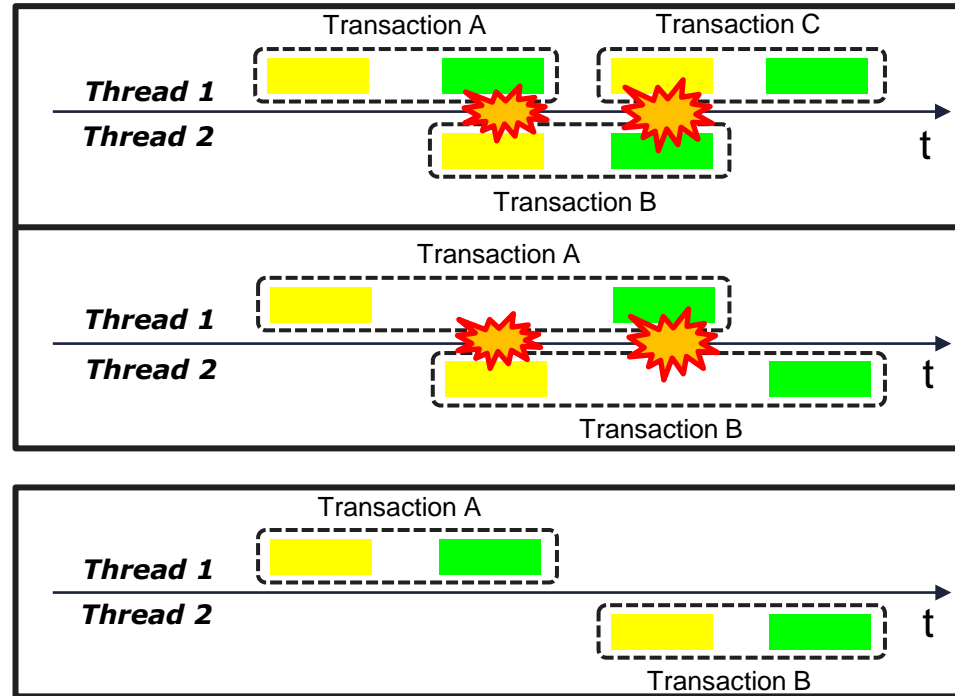
```
....
```

```
lock(&lock);;
```

```
*0x40047404 = 0x10;
```

```
unlock(&lock);;
```

```
}
```



Concurrency Bugs

lock1 != lock2

Single lock

Single lock covering each transaction

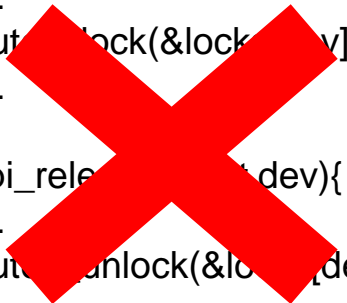
# Result: Bug Discovery

Platform	# of Bugs	# of True Positive Bugs	# of False Positive Bugs
ArduPilot	20	8	12
RaceFlight	0	0	0
RIOT	9	8	1
Contiki	0	0	0
TS100	1	1	0
grbl	0	0	0
rusEFI	6	0	6
<b>Total</b>	<b>36</b>	<b>17</b>	<b>19</b>

# Case Study

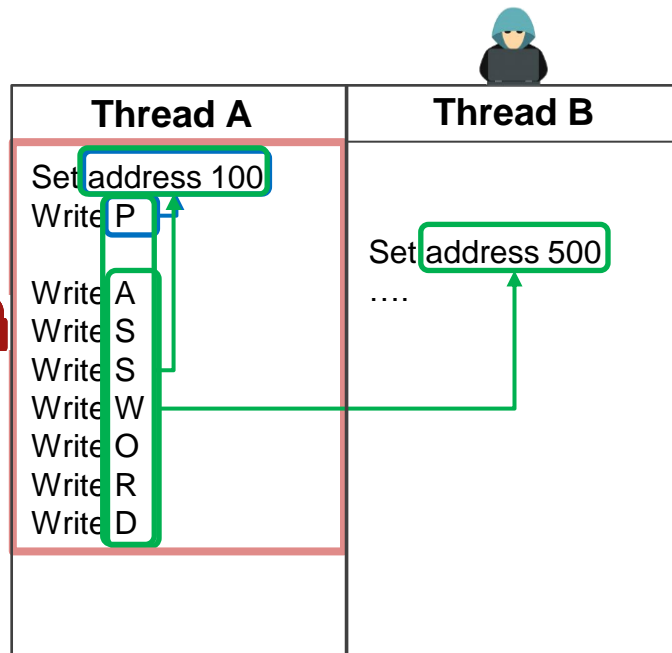
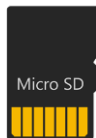
- Security-Critical Data Corruption

```
int spi_acquire(spi_t dev){  
    ....  
    mutex_lock(&lock[dev]);  
    ....  
}  
int spi_release(spi_t dev){  
    ....  
    mutex_unlock(&lock[dev]);  
    ....  
}
```



```
int sdcard_spi_write();  
int sdcard_spi_read();  
....
```

Data update  
on SD card



# Discussion

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- Limitations Inherited from Static Analysis
  - False positives and negatives caused by points-to analysis
  - Cannot handle individual interrupts
- Using Incorrectly Inferred Transaction Spans
  - Cause false positives and negatives
  - Useful to detect “transaction-level” concurrency bugs

# Conclusion

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- Concurrency bugs in embedded platforms
  - Threat to system safety and security
- PASAN: a static analysis-based concurrency bug detector
  - Targeting peripheral devices
  - Transaction-aware
    - Considering spatial and temporal atomicity
- 17 bugs on 7 real-world embedded applications

# Thank you!

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